

Proofs of transcendence

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November 2015

Objectives

- ▶ Study the gap between practical mathematics and computer-verified reasoning
- ▶ Explore structures used in various areas of mathematics
- ▶ Explore interfaces between two domains
 - ▶ Algebra: polynomials
 - ▶ Analysis: exponentiation, integration, limits
- ▶ Extend proof systems and libraries
 - ▶ Extending in the direction of multi-variate polynomials
- ▶ A long studied theme around the π number
 - ▶ Machin-like formulas, Arithmetic-geometric means, spigot

Context of work

- ▶ Subsets of complex numbers
- ▶ Polynomials in various rings
- ▶ Integration of functions with a real variable
- ▶ Multivariate polynomials
- ▶ A proof plan provided by Niven (1939)

Complex numbers

- ▶ A place where constructive mathematics take a turn
- ▶ Nijmegen experiment: C-CoRN, constructive presentation of analysis
 - ▶ No excluded middle
 - ▶ No discontinuity in functions
- ▶ Coq standard library: incursion of classical logic in type theory
 - ▶ Loose the property that proofs of existence are algorithms
 - ▶ Explore the consequences of the axioms defining real numbers
- ▶ Alternative take in Mathematical Components
 - ▶ A constructive study of real-closed fields and field extensions
 - ▶ `complex` is not a type but a type constructor

The main lemma of the proof

- ▶ If $T = c \times \prod_{i < n} (X - \alpha_i)$ has integer coefficients ($\alpha_i \neq 0$)
- ▶ And $k + \sum_{i < n} \gamma_i e^{\alpha_i} = 0$, with k, γ integers, $k \neq 0$
- ▶ Then, there exists a polynomial G with integer coefficients and degree np such that $c^{np} \sum_{i < n} \gamma_i G(\alpha_i)$ is not an integer

Using the main lemma for e

- ▶ Assume e is algebraic

$$k + \sum_{i < n} \gamma_i e^{\alpha_i} = 0 \text{ with } \alpha_i = i + 1$$

- ▶ $\prod_{i < n} (X - \alpha_i)$ trivially has integer coefficient

- ▶ For any G with integer coefficients, $\sum_{i < n} \gamma_i G(\alpha_i)$ is an integer, obviously.

Using the main lemma for π

- ▶ Assume $i\pi$ algebraic, a root of $c \times \prod_{i < n} (X - \beta_i)$

- ▶ Consider $\prod_{i < n} (1 + e^{\beta_i}) = 0$

In fact
$$\sum_{f: \{1 \dots n\} \rightarrow \{0,1\}} e^{\sum_{1 < n} f(i) \times \beta_i} = 0$$

- ▶ The α_i will be the $\sum_{1 < n} f(i) \times \beta_i \neq 0$ (n' such elements)
- ▶ $c^m \times \prod (X - \alpha_i)$ has integer coefficients by symmetry arguments
- ▶ $\gamma_i = 1$ so $c^{mn'p} \sum_i \gamma_i G(\alpha_i)$ is an integer by symmetry arguments

Symmetry arguments

- ▶ if $a_0 + a_1X + \cdots + a_nX^n = \prod_{i < n} (X - \alpha_i)$, the coefficients a_j are elementary symmetric multivariate polynomials in the α_i

$$a_j = (-1)^{n-j} \sigma_{n,j} \bar{\alpha} \quad \sigma_{n,j} = \sum_{\substack{|h|=j \\ h \subset \{1 \dots n\}}} \prod_{i \in h} X_i$$

- ▶ Example: $(X - \alpha_1)(X - \alpha_2)(X - \alpha_3) = -\alpha_1\alpha_2\alpha_3 + (\alpha_1\alpha_2 + \alpha_2\alpha_3 + \alpha_1\alpha_2)X - (\alpha_1 + \alpha_2 + \alpha_3)X^2 + X^3$
- ▶ Elementary symmetric polynomials generate all symmetric polynomial expressions
example: $X_1^2 + X_2^2 + X_3^2 = (X_1 + X_2 + X_3)^2 - 2(X_1X_2 + X_2X_3 + X_1X_3) = \sigma_{3,1}^2 - 2\sigma_{3,2}$

Proof of the main lemma

- ▶ $\int_0^1 \alpha e^{-\alpha x} P(\alpha x) dx = \sum_{i=0}^{\deg(P)} P^{(i)}(0) - e^{-\alpha} \sum_{i=0}^{\deg(P)} P^{(i)}(\alpha)$
- ▶ Name $I_P(\alpha)$ the integral, $P_d = \sum_{i=0}^{\deg(P)} P^{(i)}$
- ▶ consider $P = c^n X^{p-1} T^p$, roots of T have multiplicity p in P
 $c^{np} \sum -\gamma_i e^{\alpha_i} I_P(\alpha_i) = -c^{np} \sum \gamma_i e^{\alpha_i} P_d(0) + c^{np} \sum \gamma_i P_d(\alpha_i)$
- ▶ as p grows, the left hand side can be shown to be smaller than $(p-1)!$

Decomposing the right hand side

- ▶ Use the hypothesis $k + \sum \gamma_i e^{\alpha_i} = 0$
- ▶ Use the fact that 0 is a root with multiplicity $(p - 1)$ of P
- ▶ When $P \in \mathbb{Z}[X]$, $P^{(i)}$ has coefficients divisible by $i!$

$$-c^{np} \sum \gamma_i e^{\alpha_i} P_d(0) = kc^{np}(p-1)!T(0)^p + \sum_{i=p}^{\deg(P)} P^{(i)}(0)$$

- ▶ If p is large enough, this number is a multiple of $(p - 1)!$ but not of $p!$

Decomposing right hand side (second part)

- ▶ Use the fact that α_i are roots with multiplicity p of P

$$\sum \gamma_i P_d(\alpha_i) = \sum_i \gamma_i \sum_{j=p}^{\deg(P)} P^{(j)}(\alpha_i)$$

- ▶ Then use the fact coefficients of $P^{(j)}$ are multiples of coefficients of P and $p!$ when $p \leq j$

- ▶ Exhibit $G = \sum_{j=p}^{\deg(P)} \frac{P^{(j)}}{p!}$, G has integer coefficients

- ▶ if $c^{np} \sum_i \gamma_i G(\alpha_i)$ is an integer, then
 - ▶ the right hand side is a multiple of $(p-1)!$ and not of $p!$
 - ▶ it must be larger than $(p-1)!$, in contradiction with slide 10.

Difficulties in formalization effort

- ▶ Different definitions of complex numbers
 - ▶ Coquelicot and Math-Components each have their own hierarchy
- ▶ Exponentiation and powers
- ▶ products of sums, filters on lists
- ▶ The fundamental theorem of symmetric polynomials

Complex numbers

- ▶ Relying on the Coquelicot library
 - ▶ Coquelicot provides a 600 line-file with basic operations
 - ▶ \mathbb{C} is defined as \mathbb{R}^2
 - ▶ Properties of field, complete normed module, with a notion of derivative
- ▶ Relying on Math-components library
 - ▶ First include R (from standard library) into math-comp. hierarchy
 - ▶ Use a type-constructor: build a new field from an existing one
 - ▶ Provides the fund. th. of alg. as soon as existing field is RCF
 - ▶ Reproduce mathematical hierarchy of Coquelicot on top of math-comp

Equipment for complex numbers

- ▶ Complex integers, Complex natural numbers \mathbb{Cint} , \mathbb{Cnat}
- ▶ Generic notion of ring predicate and associated theorems

`rpred_sum`

```
: forall (V : zmodType) (S : predPredType V)
  (addS : addrPred S) (kS : keyed_pred addS)
  (I : Type) (r : seq I)
  (P : pred I) (F : I -> V),
(forall i : I, P i -> F i \in S) ->
\sum_(i <- r | P i) F i \in S
```

Exponential and power

- ▶ Important property is : $a^{(n+p)} = a^n * a^p$
- ▶ Different ways to define x^y depending on x positive or y integer or real
- ▶ real exponential: e^x defined using a power series
- ▶ complex exponential $e^{x+iy} = e^x \times (\cos y + i \sin y)$

Big operations with filters

- ▶ For π , transforming $\prod_{i < n} 1 + e^{\beta_i}$ into a sum of products of exponentials
- ▶ Example $(1 + e^{\beta_1})(1 + e^{\beta_2})(1 + e^{\beta_3})$ contains $e^{\beta_1}e^{\beta_3} = e^{\beta_1+\beta_3}$
- ▶ Discard the sums that are zero
- ▶ Lemmas already covered in the library

`bigA_distr_bigA :`

```
forall R (zero one : R) (times : Monoid.mul_law zero)
  (plus : Monoid.add_law zero times) (I J : finType)
  (F : I -> J -> R),
\big[times/one]_i \big[plus/zero]_j F i j =
\big[plus/zero]_(f : {ffun I -> J})
  \big[times/one]_i F i (f i)
```


Fundamental theorem of symmetric polynomials

- ▶ Proved for any commutative ring by P.-Y. Strub
- ▶ Stronger statement than usually found in literature: bounding degree
- ▶ Need one more refinement: preserve integer coefficients
- ▶ Rely on “morphism” between Complex numbers that are integers and integers

Symmetry arguments in more details

- ▶ We have $c \times \prod (X - \alpha_i) \in \mathbb{Z}[X]$, not $\prod (X - \alpha_i)$
- ▶ $\sigma_{n,m}(\bar{\alpha})$ is only guaranteed to be rational, not integer
Similarly $\sum G(\alpha_i)$ is only guaranteed to be rational
- ▶ But G is obtained from $T = c \times \prod (X - \alpha_i)$ by
 - ▶ Choosing an arbitrary p prime with $|c|, |k|, |T(0)|$
 - ▶ Raising to the power p
 - ▶ Multiplying by X^{p-1}
 - ▶ Computing derivatives
- ▶ Solution: multiply T by a power of c so that $\sum G(\alpha_i)$ is an integer

Lessons

- ▶ One the coming challenges is to combine libraries
 - ▶ Math-components and Coquelicot are still very close in spirit
 - ▶ Research in refactoring tools is on-going
- ▶ Navigate between types and predicates
 - ▶ `Cint` being a “ring” predicate
- ▶ Difficulties in finding the right level of abstraction
 - ▶ Very context dependent: completeness of interfaces