Formal C semantics: CompCert and the C standard

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Underspecification in C

- Unspecified behavior: two or more behaviors are allowed For example: order of evaluation in expressions
- Implementation defined behavior: like unspecified behavior, but the compiler has to document its choice For example: size and endianness of integers
- Undefined behavior: the standard imposes no requirements at all, the program is even allowed to crash For example: dereferencing a NULL or dangling pointer, signed integer overflow, ...

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- Unspecified behavior: two or more behaviors are allowed For example: order of evaluation in expressions Non-determinism
- Implementation defined behavior: like unspecified behavior, but the compiler has to document its choice For example: size and endianness of integers Parametrization
- Undefined behavior: the standard imposes no requirements at all, the program is even allowed to crash For example: dereferencing a NULL or dangling pointer, signed integer overflow, ...

No semantics/crash state

Pros and cons of underspecification

Pros for optimizing compilers:

- More optimizations are possible
- High run-time efficiency
- Easy to support multiple architectures

Cons for programmers/formal methods people:

- Portability and maintenance problems
- Hard to formally reason about

Approaches to underspecification

CompCert (Leroy *et al.*)

- Main goal: verified optimizing compiler in
- Specific choices for unspecified/impl-defined behavior For example: 32-bits ints
- Describes some undefined behavior
 For example: dereferencing NULL, integer overflow defined
- Compiler correctness proof only for programs without undefined behavior

Formalin (Krebbers & Wiedijk)

- Main goal: compiler independent separation logic in ⁹/₂
- Describes some implementation-defined behavior For example: no legacy architectures with 1s' complement
- Aims to describe all unspecified and undefined behavior

Defined behaviors in C11, Formalin and CompCert C



Defined behaviors in C11, Formalin and CompCert C



This talk: add \bigcirc to CompCert so we get Formalin \subseteq CompCert

```
void inc_array(int *p, int n) {
    int *end = p + n;
    while (p < end) (*p++)++;
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Useful:





Bizarre:

int x, y; if (&x + 1 == &y) printf("x and y are adjacent\n");

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Both undefined behavior in CompCert (1.12 and before)

Comparing with end-of-array pointers (solution)

Solution: Comparison of pointers is defined if:

Same block: both should within block bounds



Comparing with end-of-array pointers (solution)

Solution: Comparison of pointers is defined if:

Same block: both should within block bounds



Different block: both should be strictly within block bounds



Comparing with end-of-array pointers (solution)

Solution: Comparison of pointers is defined if:

Same block: both should within block bounds



Stable under compilation and gives a semantics to common programming practice with end-of-array pointers

```
struct { short x; short *r; } s1 = {10, &s.x}, s2;
unsigned char *p = &s1, *q = &s2;
unsigned char *end = p + size_of(s1);
while (p < end) *p++ = *q++;</pre>
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struct { short x; short *r; } s1 = {10, &s.x}, s2; unsigned char *p = &s1, *q = &s2; unsigned char *end = p + size_of(s1); while (p < end) *p++ = *q++;</pre>



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Byte-wise copying of objects (solution)

Solution: extend values with pointer fragment values

```
Inductive val: Type :=
  | Vundef: val
  | Vint: int -> val
  | Vlong: int64 -> val
  | Vfloat: float -> val
  | Vptr: block -> int -> val
  | Vptrfrag: block -> int -> nat -> val.
```

Subtleties:

- Dealing with arithmetic on pointer fragments
- Dealing with implicit casts (at assignments)
- More values possible, need to extend static analysis

Conclusion and future work

Semantics to useful behaviors that were previously undefined

- Comparing with end-of-array pointers
- Byte-wise pointer copy
- CompCert proofs adapted for these extensions
 - Small changes to the semantics
 - Involves proofs of many compilation passes
- Needed for cross validation of CompCert and Formalin
- Call-by-reference passing of struct values future work

Questions

Sources: http://github.com/robbertkrebbers